Biostatistics 615/815 Lecture 17: Numerical Optimization

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Annoucements

Homework

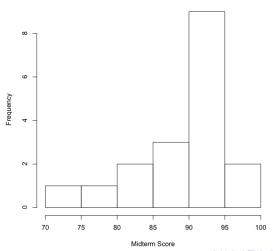
- Homework #5 will be annouced later today
- Apologies for the delay!

815 Projects

- Report the current progress to the instructore by the weekend
- Schedule a meeting with instructor by email

Midterm Score Distribution

Midterm Score Histogram (n=18)



Recap from last lecture

- Crude Monte Carlo method : calculate integration by taking averages across samples from uniform distribution
- Rejection sampling
 - 1 Define a finite rectangle
 - Sample data from uniform distribution
 - **3** Accept data if y < f(x)
 - f 4 Count how many y were hit
- Importance sampling: Reweight the probability distribution to reduce the variance in the estimation

Homework problem: integration in multivariate normal distribution

Problem

Introduction 00000

Calculate

$$\int_{x_m}^{x_M} \int_{y_m}^{y_M} f(x, y; \rho) \, dx dy$$

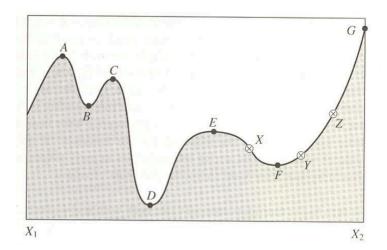
where $f(x, y; \rho)$ is pdf of bivariate normal distribution, using

- Crude Monte Carlo Method
 - Rejection sampling
 - Importance sampling

Disclaimer,

- The lecture note is very similar to Goncalo's old lecture notes
- C-specific portions are ported into C++
- The following lecture notes will be also similar.

The Minimization Problem



Specific Objectives

Finding global minimum

- The lowest possible value of the function
- Very hard problem to solve generally

Finding local minimum

- Smallest value within finite neighborhood
- Relatively easier problem

A quick detour - The root finding problem

- Consider the problem of finding zeros for f(x)
- Assume that you know
 - Point a where f(a) is positive
 - Point b where f(b) is negative
 - f(x) is continuous between a and b
- How would you proceed to find x such that f(x) = 0?

A C++ Example : definining a function object

```
#include <iostream>
class myFunc {    // a typical way to define a function object
public:
  double operator() (double x) const {
    return (x*x-1):
};
int main(int argc, char** argv) {
  myFunc foo;
  std::cout << "foo(0) = " << foo(0) << std::endl;
  std::cout << "foo(2) = " << foo(2) << std::endl;
}
```

Root Finding with C++

```
// binary-search-like root finding algorithm
double binaryZero(myFunc foo, double lo, double hi, double e) {
  for (int i=0;; ++i) {
    double d = hi - lo;
    double point = lo + d * 0.5; // find midpoint between lo and hi
    double fpoint = foo(point);  // evaluate the value of the function
    if (fpoint < 0.0) {
      d = lo - point; lo = point;
    else {
      d = point - hi; hi = point;
    // e is tolerance level (higher e makes it faster but less accruate)
    if (fabs(d) < e || fpoint == 0.0) {</pre>
      std::cout << "Iteration " << i << ", point = " << point
                << ", d = " << d << std::endl;
      return point;
```

Improvements to Root Finding

Approximation using linear interpolation

$$f^*(x) = f(a) + (x-a)\frac{f(b) - f(a)}{b-a}$$

Root Finding Strategy

• Select a new trial point such that $f^*(x) = 0$

Root Finding Using Linear Interpolation

```
double linearZero (myFunc foo, double lo, double hi, double e) {
  double flo = foo(lo); // evaluate the function at the end pointss
  double fhi = foo(hi);
  for(int i=0::++i) {
    double d = hi - lo:
    double point = lo + d * flo / (flo - fhi); //
    double fpoint = foo(point);
    if (fpoint < 0.0) {</pre>
      d = lo - point;
      lo = point:
      flo = fpoint;
    else {
      d = point - hi;
      hi = point;
      fhi = fpoint;
    if (fabs(d) < e || fpoint == 0.0) {</pre>
      std::cout << "Iteration " << i << ", point = " << point << ", d = " << d << std::endl;
      return point:
```

Performance Comparison

Finding $\sin(\mathbf{x})$ = 0 between $-\pi/4$ and $\pi/2$

```
#include <cmath>
class myFunc {
public:
    double operator() (double x) const { return sin(x); }
};
...
int main(int argc, char** argv) {
    myFunc foo;
    binaryZero(foo,0-M_PI/4,M_PI/2,1e-5);
    linearZero(foo,0-M_PI/4,M_PI/2,1e-5);
    return 0;
}
```

Experimental results

```
binaryZero(): Iteration 17, point = -2.99606e-06, d = -8.98817e-06 linearZero(): Iteration 5, point = 0, d = -4.47489e-18
```

R example of root finding

```
> uniroot( sin, c(0-pi/4,pi/2) )
$root
[1] -3.531885e-09
$f.root
[1] -3.531885e-09
$iter
[1] 4
$estim.prec
[1] 8.719466e-05
```

Summary on root finding

- Implemented two methods for root finding
 - Bisection Method : binaryZero()
 - False Position Method : linearZero()
- In the bisection method, the bracketing interval is halved at each step
- For well-behaved function, the False Position Method will converage faster, but there is no performance guarantee.

Back to the Minimization Problem

- Consider a complex function f(x) (e.g. likelihood)
- Find x which f(x) is maximum or minimum value
- Maximization and minimization are equivalent
 - Replace f(x) with -f(x)

Notes from Root Finding

- Two approaches possibly applicable to minimization problems
- Bracketing
 - Keep track of intervals containing solution
- Accuracy
 - Recognize that solution has limited precision

Notes on Accuracy - Consider the Machine Precision

- When estimating minima and bracketing intervals, floating point accuracy must be considered
- In general, if the machine precision is ϵ , the achievable accuracy is no more than $\sqrt{\epsilon}$.
- $\sqrt{\epsilon}$ comes from the second-order Taylor approximation

$$f(x) \approx f(b) + \frac{1}{2}f'(b)(x-b)^2$$

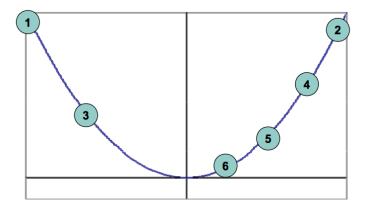
- For functions where higher order terms are important, accuracy could be even lower.
 - For example, the minimum for $f(x) = 1 + x^4$ is only estimated to about $\epsilon^{1/4}$

Outline of Minimization Strategy

- Bracket minimum
- 2 Successively tighten bracket interval

Detailed Minimization Strategy

- Find 3 points such that
 - a < b < c
 - f(b) < f(a) and f(b) < f(c)
- Then search for minimum by
 - Selecting trial point in the interval
 - Keep minimum and flanking points



Part I: Finding a Bracketing Interval

- Consider two points
 - x-values a, b
 - y-values f(a) > f(b)

Bracketing in C++

```
#define SCALE 1.618
void bracket( myFunc foo, double& a, double& b, double& c) {
  double fa = foo(a);
  double fb = foo(b);
  double fc = foo(c = b + SCALE*(b-a));
  while( fb > fc ) {
   a = b; fa = fb;
   b = c; fb = fc;
   c = b + SCALE * (b-a);
   fc = foo(c);
```

Part II: Finding Minimum After Bracketing

- Given 3 points such that
 - a < b < c
 - f(b) < f(a) and f(b) < f(c)
- How do we select new trial point?





We want to minimize the size of next search interval, which will be either from A to X or from B to C

Minimizing worst case possibility

Formulae

$$w = \frac{b-a}{c-a}$$
$$z = \frac{x-b}{c-a}$$

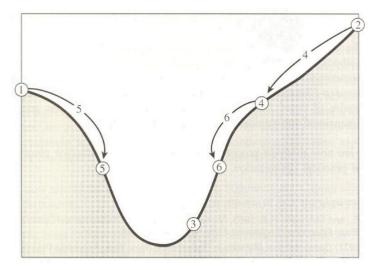
Segments will have length either 1 - w or w + z.

Optimal case

$$\frac{1-w}{z} = w+z$$

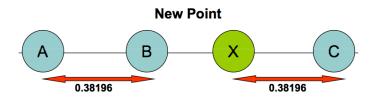
$$\frac{z}{1-w} = w$$

$$w = \frac{3-\sqrt{5}}{2} = 0.38197$$



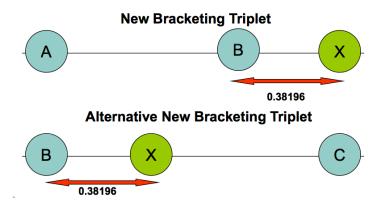


The Golden Ratio



The number 0.38196 is related to the golden mean studied by Pythagoras

The Golden Ratio



Golden Search

- Reduces bracketing by $\sim 40\%$ after function evaluation
- Performance is independent of the function that is being minimized
- In many cases, better schemes are available

Golden Step

Golden Search

```
double goldenSearch(myFunc foo, double a, double b, double c, double e) {
  int i = 0:
  double fb = foo(b);
  while ( fabs(c-a) > fabs(b*e) ) {
    double x = b + goldenStep(a, b, c);
    double fx = foo(x);
    if ( fx < fb ) {
     (x > b)? (a = b): (c = b);
     b = x; fb = fx;
   else {
     (x < b) ? (a = x) : (c = x);
    }
    ++i;
  std::cout << "i = " << i << ", b = " << b << ", f(b) = " << foo(b) << std::endl;
  return b;
}
```

A running example

Finding minimum of $f(x) = -\cos(x)$

```
class myFunc {
public:
  double operator() (double x) const {
    return 0-cos(x);
};
int main(int argc, char** argv) {
  myFunc foo;
  goldenSearch(foo,0-M PI/4,M PI/4,M PI/2,1e-5);
  return 0;
}
```

Results

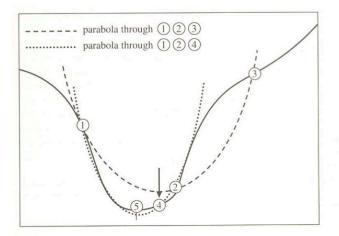
```
i = 66, b = -4.42163e-09, f(b) = -1
```

R example of minimization

```
> optimize(cos,interval=c(0-pi/4,pi/2),maximum=TRUE)
$maximum
[1] -8.648147e-07
$objective
[1] 1
```

- As with root finding, performance can improve substantially when local approximation is used
- However, a linear approximation won't do in this case.

Approximation Using Parabola



Summary

Today

- Root Finding Algorithms
 - Bisection Method : Simple but likely less efficient
 - False Position Method : More efficient for most well-behaved function
- Single-dimensional minimization
 - Golden Search

Summary

Today

- Root Finding Algorithms
 - Bisection Method : Simple but likely less efficient
 - False Position Method : More efficient for most well-behaved function
- Single-dimensional minimization
 - Golden Search

Next Lecture

- More Single-dimensional minimization
 - Brent's method
- Multidimensional optimization
 - Simplex method

