Biostatistics 615/815 Lecture 11: More Graph Algorithms Hidden Markov Models

HMMs

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Floyd-Warshall HMMs Forward-backward Viterbi

Annoucement

Introduction

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- Homework 3 is announced, due next Tuesday.
- Try to install boost library by today and ask technical questions tomorrow if there is any



Recap: boost library

Introduction

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```
#include <iostream>
#include <boost/tokenizer.hpp>
#include <string>
using namespace std;
using namespace boost;
int main(int argc, char** argv) {
 // default delimiters are spaces and punctuations
  string s1 = "Hello, boost library";
 tokenizer<> tok1(s1);
 for(tokenizer<>::iterator i=tok1.begin(); i != tok1.end(); ++i) {
    cout << *i << endl:
 // you can parse csv-like format
  string s2 = "Field 1,\"putting quotes around fields, allows commas\",Field 3";
 tokenizer<escaped list separator<char> > tok2(s2);
 for(tokenizer<escaped list separator<char> >::iterator i=tok2.begin();
       i != tok2.end(); ++i) {
    cout << *i << endl:
 return 0;
}
```

Recap: Dijkstra's algorithm

Algorithm DIJKSTRA

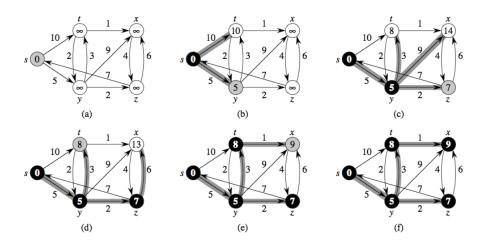
Introduction

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```
Data: G: graph, w: weight, s: source
Result: Each vertex contains the optimal weight from s
INITIALIZESINGLESOURCE(G,s);
S = \emptyset:
Q = G.V:
while Q \neq \emptyset do
    u = \text{EXTRACTMIN}(Q);
    S = S \cup \{u\};
   for v \in G.Adi[u] do
       Relax(u, v, w);
    end
```

end

Recap: Illustration of DIJKSTRA's algorithm





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Calculating all-pair shortest-path weights

A dynamic programming formulation

Let $d_{ij}^{(k)}$ be the weight of shortest path from vertex i to j, for which intermediate vertices are in the set $\{1, 2, \dots, k\}$.

$$d_{ij}^{(k)} = \begin{cases} w_{ij} & k = 0\\ \min(d_{ij}^{(k-1)}, d_{ik}^{(k-1)} + d_{ki}^{(k-1)}) & k > 0 \end{cases}$$

Algorithm FLOYDWARSHALL

```
Data: W: n \times n weight matrix
D^{(0)} = W;
for k = 1 to n do
     for i = 1 to n do
          for j = 1 to n do
               d_{ii}^{(k)} = \min(d_{ii}^{(k-1)}, d_{ik}^{(k-1)} + d_{ki}^{(k-1)});
          end
     end
end
return D^{(n)}:
```

Summary: shortest path finding algorithms

Dijkstra's algorithm

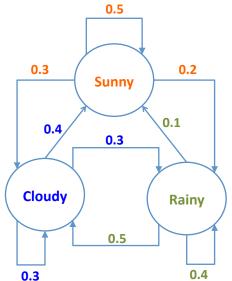
- $\Theta(|V| \log |V| + |E|)$ dynamic programming algorithm.
- Compute optimal path from a single source to each node
- Track optimal path from the closest node from the source, and expand to adjacent node

Floyd-Warshall algorithm

- \bullet $\Theta(|\mathit{V}|^3)$ All-pair shortest path finding algorithms with non-negative weights
- \bullet Use the fact that the maximum length of each possible optimal path is $\mid V \mid.$
- \bullet For each possible pairs of sources and destinations, iteratively update optimal distance matrix $\mid V \mid$ times.

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Markov Process : An example





$$\pi = \begin{pmatrix} \Pr(q_1 = S_1 = \mathsf{Sunny}) \\ \Pr(q_1 = S_2 = \mathsf{Cloudy}) \\ \Pr(q_1 = S_3 = \mathsf{Rainy}) \end{pmatrix} = \begin{pmatrix} 0.7 \\ 0.2 \\ 0.1 \end{pmatrix}$$

$$A_{ij} = \Pr(q_{t+1} = S_i | q_t = S_j)$$

$$A = \begin{pmatrix} 0.5 & 0.4 & 0.1 \\ 0.3 & 0.3 & 0.5 \\ 0.2 & 0.3 & 0.4 \end{pmatrix}$$

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Stationary distribution

$$\mathbf{p} = A\mathbf{p} p = (0.346, 0.359, 0.295)^T$$

Markov process is only dependent on the previous state

If it rains today, what is the chance of rain on the day after tomorrow?

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If it rains today, what is the chance of rain on the day after tomorrow?

$$\Pr(q_3 = S_3 | q_1 = S_3) = \begin{bmatrix} A^2 & 0 \\ 0 \\ 1 \end{bmatrix}_3 = 0.33$$

If it has rained for the past three days, what is the chance of rain on the day after tomorrow?

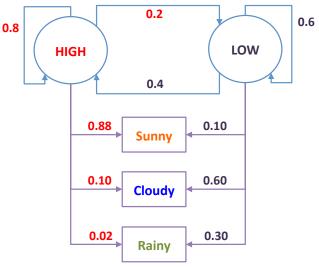
$$\Pr(q_5 = S_3 | q_1 = q_2 = q_3 = S_3) = \Pr(q_5 = S_3 | q_3 = S_3) = 0.33$$

Hidden Markov Models (HMMs)

- A Markov model where actual state is unobserved
 - Transition between states are probablistically modeled just like the Markov process
- Typically there are observable outputs associated with hidden states
 - The probability distribution of observable outputs given an hidden states can be obtained.

troduction Floyd-Warshall **HMMs** Forward-backward Viterbi

An example of HMM



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States
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HMMs 0000000000000

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Initial States $\pi_i = \Pr(q_1 = S_i), \ \pi = \{0.7, 0.3\}$
Transition $A_{ii} = \Pr(q_{t+1} = S_i | q_t = S_i)$

$$A = \left(\begin{array}{cc} 0.8 & 0.4 \\ 0.2 & 0.6 \end{array}\right)$$

Emission
$$B_{ij} = b_{q_t}(o_t) = b_{S_j}(O_i) = \Pr(o_t = O_i | q_t = S_j)$$

$$B = \left(\begin{array}{cc} 0.88 & 0.10\\ 0.10 & 0.60\\ 0.02 & 0.30 \end{array}\right)$$



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- What is the chance of rain in the day 3?
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- What is the chance of rain in the day 3, if it rained in the day 1 and day 2?
- If the observation was (SUNNY,SUNNY,CLOUDY,RAINY,RAINY) from day 1 through day 5, what would be the mostly likely sequence of states?
- Can we infer the HMM paremeters if we have a large number of observations?



Unconditional marginal probabilities

What is the chance of rain in the day 4?

$$\mathbf{f}(\mathbf{q}_3) = \begin{pmatrix} \Pr(q_4 = S_1) \\ \Pr(q_4 = S_2) \end{pmatrix} = A^3 \pi = \begin{pmatrix} 0.669 \\ 0.331 \end{pmatrix}$$

$$\mathbf{g}(o_4) = \begin{pmatrix} \Pr(o_4 = O_1) \\ \Pr(o_4 = O_2) \\ \Pr(o_4 = O_3) \end{pmatrix} = B\mathbf{f}(\mathbf{q}_4) = \begin{pmatrix} 0.621 \\ 0.266 \\ 0.233 \end{pmatrix}$$

The chance of rain in day 3 is 23.3%

Calculating conditional probabilities

What is the chance of rain in the day 2 if it rains in the day 1?

$$\Pr(o_2 = O_3 | o_1 = O_3) = \Pr(o_2 = O_3 | q_2 = S_1) \Pr(q_2 = S_1 | o_2 = O_3) + \Pr(o_2 = O_3 | q_2 = S_2) \Pr(q_2 = S_2 | o_2 = O_3)$$
...

This is already quite complicated!



Organizing the likelihood

- Let $\lambda = (A, B, \pi)$
- For a sequence of observation $\mathbf{o} = \{o_1, \dots, o_t\},\$

$$\begin{split} & \Pr(\mathbf{o}|\lambda) &= \sum_{\mathbf{q}} \Pr(\mathbf{o}|\mathbf{q},\lambda) \Pr(\mathbf{q}|\lambda) \\ & \Pr(\mathbf{o}|\mathbf{q},\lambda) &= \prod_{i=1}^t \Pr(o_i|q_i,\lambda) = \prod_{i=1}^t b_{q_i}(o_i) \\ & \Pr(\mathbf{q}|\lambda) &= \pi_{q_1} \sum_{i=2}^t a_{q_iq_{i-1}} \\ & \Pr(\mathbf{o}|\lambda) &= \sum_{\mathbf{q}} \pi_{q_1} b_{q_1}(o_{q_1}) \prod_{i=2}^t a_{q_iq_{i-1}} b_{q_i}(o_{q_i}) \end{split}$$

Naive computation of the likelihood

$$\Pr(\mathbf{o}|\lambda) = \sum_{\mathbf{q}} \pi_{q_1} b_{q_1}(o_{q_1}) \prod_{i=2}^t a_{q_i q_{i-1}} b_{q_i}(o_{q_i})$$

- Number of possible $q=2^t$ are exponentially growing with the number of observations
- Computational would be infeasible for large number of observations
- Algorithmic solution required for efficient computation.

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Forward-backward algorithm

• Define forward probability $\alpha_t(i)$ as

$$\alpha_t(i) = \Pr(o_1, \cdots, o_t, q_t = S_i | \lambda)$$

- $\alpha_t(i)$ can be efficiently computed using dynamic programming
 - $\alpha_1(i) = \pi_i b_i(o_1)$
 - $\alpha_t(i) = \sum_{j=1}^n \alpha_{t-1}(j) a_{ij} b_i(o_t)$
 - $\Pr(\mathbf{o}|\lambda) = \sum_{i=1}^{n} \alpha_t(i)$
- Time complexity is $\Theta(n^2t)$.

Forward-backward algorithm (cont'd)

• Backward probability $\beta_t(i)$ as

$$\beta_t(i) = \Pr(o_{t+1}, \cdots, o_T | q_t = S_i, \lambda)$$

- ullet $eta_t(i)$ can also be efficiently computed using dynamic programming
 - $\beta_T(i) = 1$
 - $\beta_t(i) = \sum_{j=1}^n a_{ji} b_j(o_{t+1}) \beta_{i+1}(j)$
- Time complexity is $\Theta(n^2(T-t))$.

Forward-backward algorithm (cont'd)

 We can infer the conditional probability of each state given observations by

$$\gamma_t(i) = \Pr(q_t = S_i | \mathbf{o}, \lambda)$$

$$= \frac{\Pr(\mathbf{o}, q_t = S_i | \lambda)}{\sum_{j=1}^n \Pr(\mathbf{o}, q_t = S_j | \lambda)}$$

$$= \frac{\alpha_t(i)\beta_t(i)}{\sum_{j=1}^n \alpha_t(j)\beta_t(j)}$$

• Time complexity is $\Theta(n^2 T)$.

Viterbi ●○

Viterbi algorithm

- Finding the most likely trajactory of states given a series of observations
- Want to compute

$$rg \max_{\mathbf{q}} \Pr(\mathbf{q}|\mathbf{o},\lambda)$$

• Define $\delta_t(i)$ as

$$\delta_t(i) = \max_{\mathbf{q}} \Pr(\mathbf{q}, \mathbf{o} | \lambda)$$

• Use dynamic programming algorithm to find the 'shortest' path

Viterbi 0

Viterbi algorithm (cont'd)

Initialization $\delta_1(i) = \pi b_i(o_1)$ for $1 \le i \le n$.

Maintenance $\delta_t(i) = \max_i \delta_{t-1}(i) a_{ii} b_i(o_t)$

Termination Max likelihood is $\max_i \delta_T(i)$

Reconstruction How to reconstruct the optimal path?