Biostatistics 615/815 Lecture 13: Programming with Matrix

Hyun Min Kang

February 17th, 2011



 Introduction
 Power
 Matrix
 Matrix Computation
 Linear System
 Least square
 Summary

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Annoucements

Homework #3

- Homework 3 is due today
- If you're using Visual C++ and still have problems in using boost library, you can ask for another extension

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Annoucements

Homework #3

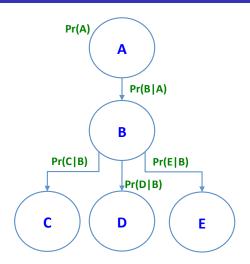
- Homework 3 is due today
- If you're using Visual C++ and still have problems in using boost library, you can ask for another extension

Homework #4

- Homework 4 is out
- Floyd-Warshall algorithm
 - Note that some key information was not covered in the class.
- Fair/biased coint HMM



Last lecture - Conditional independence in graphical models



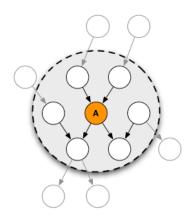
• Pr(A, C, D, E|B) = Pr(A|B) Pr(C|B) Pr(D|B) Pr(E|B)

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Markov Blanket



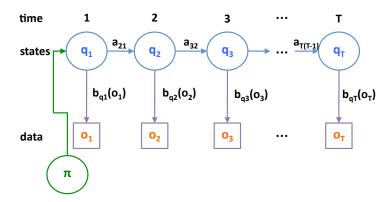
- If conditioned on the variables in the gray area (variables with direct dependency), A is independent of all the other nodes.
- $A \perp (U-A-\pi_A)|\pi_A$

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Hidden Markov Models



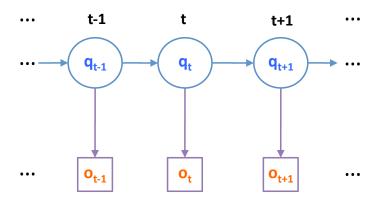


Conditional dependency in forward-backward algorithms

• Forward : $(q_t, o_t) \perp \mathbf{o}_t^- | \mathbf{q}_{t-1}$.

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• Backward : $o_{t+1} \perp \mathbf{o}_{t+1}^+ | \mathbf{q}_{t+1}$.

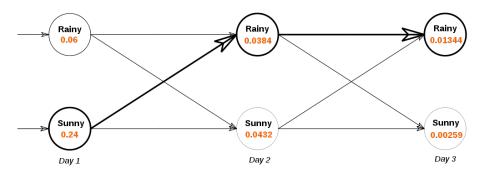


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Viterbi algorithm - example

- When observations were (walk, shop, clean)
- Similar to Dijkstra's or Manhattan tourist algorithm



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Today's lecture

- Calculating Power
- Linear algebra 101
- Using Eigen library for linear algebra
- Implementing a simple linear regression



Calculating power

Problem

- Computing a^n , where $a \in \mathbb{R}$ and $n \in \mathbb{N}$.
- How many multiplications would be needed?

Function slowPower()

```
double slowPower(double a, int n) {
  double x = a;
  for(int i=1; i < n; ++i) {
    x *= a;
  }
  return x;
}</pre>
```



More efficient computation of power

Function fastPower()

```
double fastPower(double a, int n) {
  if ( n == 1 ) {
    return a;
  else {
    double x = fastPower(a,n/2);
    if ( n % 2 == 0 ) {
      return x * x;
    else {
      return x * x * a;
```

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Computational time

```
main()
int main(int argc, char** argv) {
  double a = 1.0000001:
  int n = 10000000000;
  clock t t1 = clock();
  double x = slowPower(a,n);
  clock t t2 = clock();
  double y = fastPower(a,n);
  clock t t3 = clock():
  std::cout << "slowPower ans = " << x << ", sec = "
            << (double)(t2-t1)/CLOCKS PER SEC << std::endl;
  std::cout << "fastPower ans = " << y << ", sec = "
            << (double)(t3-t2)/CLOCKS PER SEC << std::endl:
```

Running examples

```
slowPower ans = 2.6881e+43, sec = 1.88659
fastPower ans = 2.6881e+43, sec = 3e-06
```

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Summary - fastPower()

- $\Theta(\log n)$ complexity compared to $\Theta(n)$ complexity of slowPower()
- Similar to binary search vs linear search
- Good example to illustrate how the efficiency of numerical computation could change by clever algorithms



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Programming with Matrix

Why Matrix matters?

- Many statistical models can be well represented as matrix operations
 - Linear regression
 - Logistic regression
 - Mixed models
- Efficient matrix computation can make difference in the practicality of a statistical method
- Understanding C++ implementation of matrix operation can expedite the efficiency by orders of magnitude



Ways to Matrix programmming

- Implementing Matrix libraries on your own
 - Implementation can well fit to specific need
 - Need to pay for implementation overhead
 - Computational efficiency may not be excellent for large matrices



Ways to Matrix programmming

- Implementing Matrix libraries on your own
 - Implementation can well fit to specific need
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- Using BLAS/LAPACK library
 - Low-level Fortran/C API
 - ATLAS implementation for gcc, MKL library for intel compiler (with multithread support)
 - · Used in many statistical packages including R
 - Not user-friendly interface use.
 - boost supports C++ interface for BLAS



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Ways to Matrix programmming

- Implementing Matrix libraries on your own
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- Using BLAS/LAPACK library
 - Low-level Fortran/C API
 - ATLAS implementation for gcc, MKL library for intel compiler (with multithread support)
 - Used in many statistical packages including R
 - Not user-friendly interface use.
 - boost supports C++ interface for BLAS
- Using a third-party library, Eigen package
 - A convenient C++ interface
 - Reasonably fast performance
 - Supports most functions BLAS/LAPACK provides



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Using a third party library

Downloading and installing Eigen package

- Download at http://eigen.tuxfamily.org/index.php?title=3.0_beta
- To install just uncompress it



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Using a third party library

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Using Eigen package

- Add -I DOWNLOADED_PATH/eigen option when compile
- No need to install separate library. Including header files is sufficient

Example usages of Eigen library

```
#include <iostream>
#include <Eigen/Dense> // For non-sparse matrix
using namespace Eigen; // avoid using Eigen::
int main()
                      // 2x2 matrix type is defined for convenience
 Matrix2d a:
 a << 1, 2,
       3, 4;
 MatrixXd b(2.2): // but you can define the type from arbitrary-size matrix
 b << 2, 3,
      1, 4;
  std::cout << "a + b = \n" << a + b << std::endl: // matrix addition
  std::cout << "a - b =\n" << a - b << std::endl; // matrix subtraction
  std::cout << "Doing a += b:" << std::endl:
 a += b:
  std::cout << "Now a =\n" << a << std::endl;
 Vector3d v(1,2,3):
                                                   // vector operations
 Vector3d w(1,0,0);
  std::cout << "-v + w - v = n" << -v + w - v << std::endl:
}
```

More examples

```
#include <iostream>
#include <Eigen/Dense>
using namespace Eigen;
int main()
{
                              // 2*2 matrix
  Matrix2d mat:
  mat << 1, 2,
         3, 4:
  Vector2d u(-1,1), v(2,0); // 2D vector
  std::cout << "Here is mat*mat:\n" << mat*mat << std::endl;</pre>
  std::cout << "Here is mat*u:\n" << mat*u << std::endl:
  std::cout << "Here is u^T*mat:\n" << u.transpose()*mat << std::endl;</pre>
  std::cout << "Here is u^T*v:\n" << u.transpose()*v << std::endl;
  std::cout << "Here is u*v^T:\n" << u*v.transpose() << std::endl;</pre>
  std::cout << "Let's multiply mat by itself" << std::endl;</pre>
  mat = mat*mat:
  std::cout << "Now mat is mat:\n" << mat << std::endl;</pre>
```

Time complexity of matrix computation

Square matrix multiplication / inversion

- Naive algorithm : $O(n^3)$
- Strassen algorithm : $O(n^{2.807})$
- Coppersmith-Winograd algorithm : $O(n^{2.376})$ (with very large constant factor)

Determinant

- Laplace expansion : O(n!)
- LU decomposition : $O(n^3)$
- Bareiss algorithm : $O(n^3)$
- Fast matrix multiplication algorithm : $O(n^{2.376})$



Computational considerations in matrix operations

Avoiding expensive computation

• Computation of $\mathbf{u}'AB\mathbf{v}$



Computational considerations in matrix operations

Avoiding expensive computation

- Computation of $\mathbf{u}'AB\mathbf{v}$
- If the order is $(((\mathbf{u}'(AB))\mathbf{v})$
 - $O(n^3) + O(n^2) + O(n)$ operations
 - $O(n^2)$ overall



Computational considerations in matrix operations

Avoiding expensive computation

- Computation of $\mathbf{u}'AB\mathbf{v}$
- If the order is $(((\mathbf{u}'(AB))\mathbf{v})$
 - $O(n^3) + O(n^2) + O(n)$ operations
 - $O(n^2)$ overall
- If the order is $(((\mathbf{u}'A)B)\mathbf{v})$
 - Two $O(n^2)$ operations and one O(n) operation
 - O(n²) overall



Quadratic multiplication

Same time complexity, but one is slightly more efficient

- Computing $\mathbf{x}'A\mathbf{y}$.
- $O(n^2) + O(n)$ if ordered as $(\mathbf{x}'A)\mathbf{y}$.
- \bullet Can be simplified as $\sum_i \sum_j x_i A_{ij} y_j$

A symmetric case

- Computing $\mathbf{x}'A\mathbf{x}$ where A=LL'
- $\mathbf{u} = L'\mathbf{x}$ can be computed more efficiently than $A\mathbf{x}$.
- $\mathbf{x}' A \mathbf{x} = \mathbf{u}' \mathbf{u}$



Solving linear systems

Problem

Find \mathbf{x} that satisfies $A\mathbf{x} = \mathbf{b}$

A simplest approach

- Calculate A^{-1} , and $\mathbf{x} = A^{-1}\mathbf{b}$
- Time complexity is $O(n^3) + O(n^2)$
- A has to be invertible
- Potential issue of numerical instability



Using matrix decomposition to solve linear systems

LU decomposition

- A = LU, making $U\mathbf{x} = \mathbf{L}^{-1}\mathbf{b}$
- A needs to be square and invertible.
- Fewer additions and multiplications
- Precision problems may occur

QR decomposition

- A = QR where A is $m \times n$ matrix
- Q is orthogonal matrix, Q'Q = I.
- R is $m \times n$ upper-triangular matrix, $R\mathbf{x} = Q'\mathbf{b}$.



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Cholesky decomposition

- A is a square, symmetric, and positive definite matrix.
- A = U'U is a special case of LU decomposition
- Computationally efficient and accurate



Solving least square

Solving via inverse

- Most straightforward strategy
- $\mathbf{y} = X\beta + \epsilon$, \mathbf{y} is $n \times 1$, X is $n \times p$.
- $\beta = (X'X)^{-1}X'y$.
- Computational complexity is $O(np^2) + O(np) + O(p^3)$.
- The computation may become unstable if X'X is singular
- Need to make sure that rank(X) = p.



Singular value decomposition

Definition

A $m \times n (m \ge n)$ matrix A can be represented as $A = UDV^T$ such that

- U is $m \times n$ matrix with orthogonal columns $(U'U = I_n)$
- ullet D is n imes n diagnonal matrix with non-negative entries
- V^T is $n \times n$ matrix with orthogonal matrix ($V'V = VV' = I_n$).

Computational complexity

- $4m^2n + 8mn^2 + 9m^3$ for computing U, V, and D.
- $4mn^2 + 8n^3$ for computing V and D only.
- The algorithm is numerically very stable



Stable inferecne of least square using SVD

$$X = UDV'$$

$$\beta = (X'X)^{-1}X'\mathbf{y}$$

$$= (VDU'UDV')^{-1}VDU'\mathbf{y}$$

$$= (VD^2V')^{-1}VDU'\mathbf{y}$$

$$= VD^{-2}V'VDU'\mathbf{y}$$

$$= VD^{-1}U'\mathbf{y}$$



Stable inferecne of least square using SVD

```
#include <iostream>
#include <Eigen/Dense>
using namespace std;
#using namespace Eigen;
int main()
{
   MatrixXf A = MatrixXf::Random(3, 2):
   cout << "Here is the matrix A:\n" << A << endl;</pre>
   VectorXf b = VectorXf::Random(3):
   cout << "Here is the right hand side b:\n" << b << endl:</pre>
   cout << "The least-squares solution is:\n"</pre>
        << A.jacobiSvd(ComputeThinU | ComputeThinV).solve(b) << endl;
```

ntroduction Power Matrix Matrix Computation Linear System Least square **Summary**

Summary

- Calculating Power
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- Implementing a simple linear regression

